





Maintenance of Reliable Distributed Applications with Open Source Middleware: Fifteen years later

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- Introduction
- Background
- CORBA Overview
- Maintenance Experiences
- Patching the middleware
- Improving performance
- Conclusions





- Maintenance of critical applications
 - Long life cycle (several years)
 - An important choice
 - Hardware
 - Software
 - No real data available for new technologies
 - Technologies are constantly evolving
- At the beginning of the previous decade
 - Middleware as a promising technology
 - CORBA, Java-RMI, DCOM
 - CORBA to solve all heterogeneity problems
 - Different languages, operating systems, platforms
 - RT-CORBA, FT-CORBA and Minimum CORBA





Advantages

- Some complexity is hidden by the middleware
 - Low level details
 - Network complexity
- Better interoperability
- Focus on functionality

Disadvantages

- The middleware has the "control" of the application
- We depend on the middleware and its updates
- How can this affect the maintenance? (specially how it affects software reliability)
 - Changes in operating system, languages, patches, ...

Introduction



- Current situation
 - CORBA is not as popular as expected
 - Several reasons
 - Internet (Web services, REST, ...)
 - New languages and platforms (e.g. C#, .NET)
 - Not widely accepted by the industry and users (learning curve)
 - A niche in several sectors: telecommunications, defense, simulation, ...
- But, some projects have already started...
 - And they are being exploited
 - **2**4/7
 - Two possibilities
 - To change technologies
 - But to change is not always a possibility
 - Maintenance activities

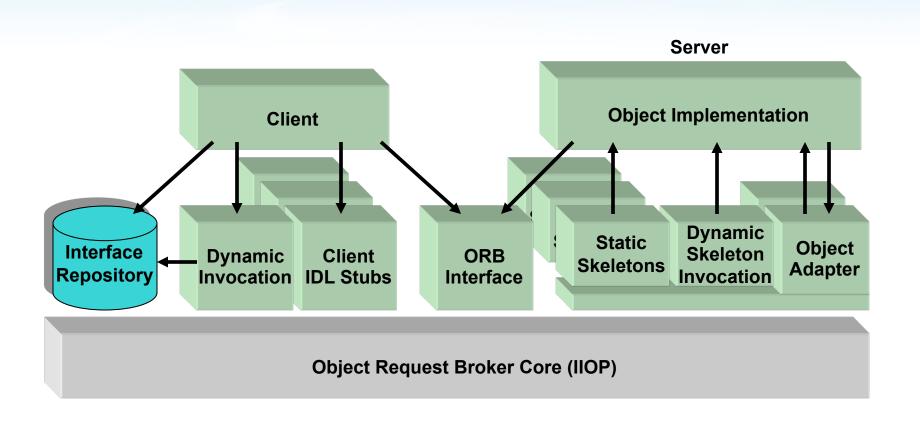




- Development of applications and communications
 - UML, software components and CORBA
- Joint projects with several companies
- UM-RTCOM was presented
 - RT-CORBA based
 - Component model
 - Real-time support
 - Higher abstraction level
 - Some experiences
 - Distributed simulators with soft real-time constraints
 - SMEPP European Project
- Two principally used CORBA implementations
 - TAO (C++)
 - JacORB (Java)



CORBA Overview General CORBA Architecture





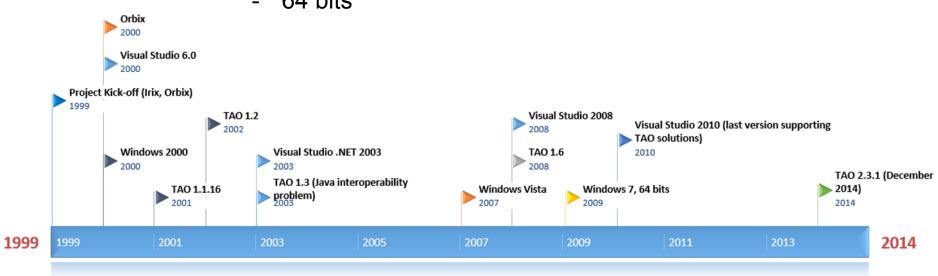


- Last CORBA version
 - 3.3, November 2012
- Many CORBA implementations are currently active (free and commercial)
 - TAO
 - JacORB
 - Orbix, Orbacus, ...
 - Java SDK implementation
- New languages have been added to CORBA
 - Python, Ruby

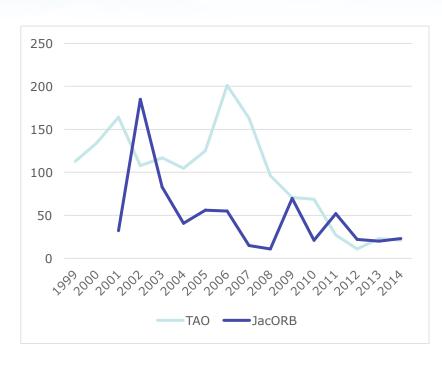


Main milestones:

- Windows Vista
- C#/.NET
- 64 bits



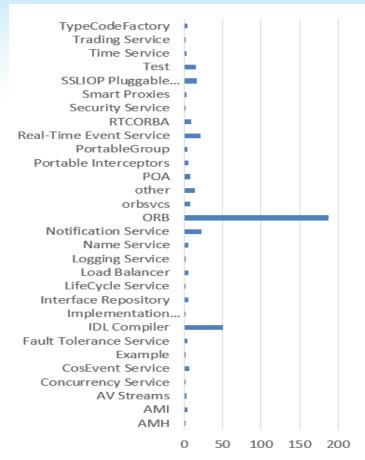




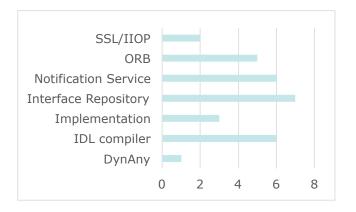
Number of bugs fixed 1999-2014

- Decreasing activity over the last few years
- A significant number of bugs remain
- Latest TAO version (2.3.2)
 - May 2015
- Latest JacORB version (3.6.1)
 - May 2015





TAO active bugs (end 2014)



JacORB active bugs (end 2014)



TAO and JacORB philosophy:

- Bugzilla: bug tracking system
- Users can contribute with solutions and improvements

Alternatively: commercial support (OCI, Remedy IT)

But, when a change is required or a fault is detected...

	Files	Lines	Statements	Class Defs
ACE+TAO	22,915	1,859,251	616,530	8,078
JacORB	2,138	300,129	115,547	2,532

ACE+TAO 6.3.1 and JacORB metrics



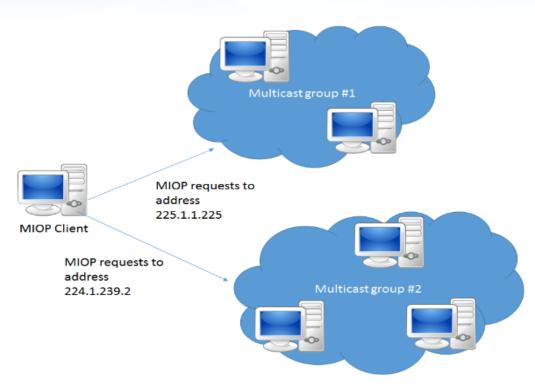
- What about testing?
 - Automatic testing is supported
 - More than 1,000 tests are included with ACE+TAO
 - Different categories
 - ACE
 - TAO
 - ORB services
 - But this is not enough in critical systems



- Improvements over time
 - Service orientation is the new trend
 - Heterogeneity is a fact
 - Communication very well encapsulated
 - Dynamism
 - Performance
 - Reliability
 - Scalability
- Different kind of changes
 - Patching the middleware
 - Improving performance
- TAO is very extensible and configurable
 - Several design patterns
 - Many configuration options



- Improvement performance (multicast)
 - MIOP/UIPMC protocols



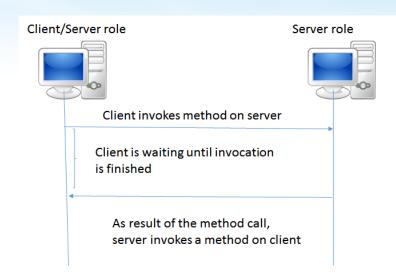
dynamic UIPMC_Factory
Service_Object *
TAO_PortableGroup:_make
_TAO_UIPMC_Protocol_Fa
ctory()

"-ORBListenOnAll 0

-ORBListenerInterfaces 224.1.239.2=192.168.20.13 5"



- Improvement scalability (concurrency control)
 - On the server side:
 - ORB controlled model: requests are attended to in the order specified by the ORB.
 - Single thread model.
 - On the client side:
 - Leader-follower: while waiting, client threads can be reused to process other requests.
 - Reactive: thread provided by the TAO Reactor
 - Blocking: the client is blocked until the connection finishes.



Deadlock problems with nested upcalls



- Customer request: "When using IPv6 multicast addresses in a PC equipped with multiple network cards, TAO ignores user preferences and it always uses the first network interface. In addition, we have a strange message"
- TAO configuration says: user can select in which network interface requests are attended.

"for a machine with two network cards identified by the ip addresses 192.168.1.10 and 192.168.1.20, you can use the single directive - ORBListenerInterfaces 239.255.*=*10,224.255.*=*20..."



- How the problem was detected...
 - TAO logs showed user preferences, but these logs were false!
 - Netstat -g shows which network interface was really used
- Where we should look for... several candidates
 - ACE sockets
 - TAO protocols
 - Configuration
 - Experiences with TAO helps a little, but a slow process



- What we found...
 - A "strange" error message related to IPv6 addresses
 - "ACE_INET_Addr::get_ip_address: address is a IPv6 address not IPv4"
 - Reason: call to addr.get_ip_address (), method only available for IPv4
 - We detected unused code. In fact, two operations can be removed from ACE+TAO
 - ACE_UINT32 uint_ip_addr (void) const;
 - void uint_ip_addr (ACE_UINT32 ip_addr);
 - Second step: when this code was removed, we obtained a "core". Reason:
 - Buffer sizes when using IPv6 string addresses
 - A constant (MAX_ADDR_LENGTH) was defined with size 32, which is not valid for IPv6 address. An example:
 - 2001:0db8:85a3:0042:1000:8a2e:192.168.158.190



- What else?
 - When these problems were solved, the user configuration continued to be ignored
 - Reason: TAO can use the following configuration syntax:
 ./server -ORBId ORB_LAN_1 ORBAllowZIOPNoServerPolicies 1 -ORBListenEndpoints
 iiop://[2001:db8:0:f101::1] -ORBPreferredInterfaces
 *=eth2 -ORBEnforcePreferredInterfaces 1 -ORBDebugLevel
 - Network interface is selected using "=interface_name"
 - Finally, we found...



- ACE_OS::if_nametoindex calls the standard function if_nametoindex(), that returns the index of the network interface corresponding to the name ifname
- TAO was using "if=ethX" as name of the network interface!!!
- When this code was changed, Voilà! It worked



- Some improvements can be applied to the code in order to get a better performance
- The following is ongoing work
- Customer request: CORBA sequences performance does not seem to be very efficient. The memset function is intensively used.



```
IDL CORBA test interface:
module MyModule {
 const long MSG_MAX_DATA_SIZE=52428800;
 typedef sequence<string, MSG_MAX_DATA_SIZE>
            ByteSequence;
 valuetype mivaluetype {
      public ByteSequence data;
 };
 interface Basic {
  void receivevt(in any vt);
  void shutdown();
```



```
Client fragment:
OBV_MyModule::mivaluetype msg;
CORBA:: Any miany;
miany <<= &msg;
for(int i=0;i<100; i++) {
     tst->receivevt(miany);
Servant fragment
void basic_i::receivevt (const CORBA::Any &vt) {
  MyModule::mivaluetype *msg;
   vt >>= msg;
```



- What valgrind and callgrind tools say:

 3,145,730,764 ???:CORBA::string_free(char*) [/home/dgarrido/ACE_wrappers/TAO/tao/libTAO.so.2.3.2]

 2,621,440,370 ???:OBV_mimodulo::mivaluetype::~mivaluetype()'2 [/home/dgarrido/src/secuencias_ms/server]

 393,222,493 ???:__GI_memset [/lib64/libc-2.12.so]
- Reason: a myriad of C++ templates defining other templates, using templates, Very hard to read! Finally, a memset operation is performed
 - What will happen with this? The memset function is expected to be very efficient



Problem: when the sequence is extracted in the servant, the memset function is called using the **MAXIMUM** size of the sequence without taking into account the actual sequence length (even when the actual length is zero!)

And this is a very costly operation when repeated several times per second with a huge buffer

Conclusions



- Fifteen years of experiences using open-source middleware
- Not a bad experience
- The importance of the choice
 - Changes in operating system, languages, platform
- New trends in middleware technology
 - Data Distribution Service (OMG)
 - Publish-subscribe mechanism
 - Not very novel (from 2003!), but the life cycle of critical applications seem different from other kinds of applications